# Oblivious Sketching of High-Degree Polynomial Kernels 

Thomas Ahle ITU

Michael Kapralov EPFL

Jakob Knudsen<br>Univ. of Copenhagen

Rasmus Pagh
ITU

Ameya Velingker David Woodruff<br>Google AI<br>CMU<br>Amir Zandieh EPFL

# ㄷ D ㄷ Carnegie Mellon University 

IT-UNIVERSITETET I KØBENHAVN

## Oblivious Sketching of High-Degree Polynomial Kernels

Thomas Ahle<br>ITU

Mich nger Google AI

David Woodruff CMU

Jakob Knudsen<br>Univ. of Copenhagen

## Rasmus Pagh <br> ITU

# Oblivious Sketching of High-Degree Polynomial Kernels 

## Thomas Ahle ITU

Michael Kapralov EPFL

Jakob Knudsen
Univ. of Conenhagen

Rasmus Pagh
ITU

Ameya Velingker David Woodruff Amir Zandieh Google AI<br>CMU

Carnegie Mellon University
Google AI

## Kernel Methods

- Widely used in kernel-based learning, statistics, and control
- Classical machine learning tool with real-world applications



## Real-World Applications of Kernel Methods

- Hyperparameter tuning of deep neural networks: e.g. Google Vizier
- Multi-Armed Bandit Optimization [Srinivas, Krause, Kakade, Seeger' 09]
- Neural Tangent Kernel: The evolution of a neural network during training can be described by kernel methods [Jacot, Gabriel, Hongler'18]


## Kernel Methods

- Learn a nonlinear function $f: \mathcal{R}^{d} \rightarrow \mathcal{R}$ from noisy samples

$$
\gamma_{i}=f\left(x_{i}\right)+\epsilon_{i} \text { for } i=1,2, \ldots n
$$

- $\epsilon_{i}$ are iid Gaussian noise with zero mean and variance $\lambda$
- Kernel Ridge Regression is a simple and yet powerful solution


## Kernel Methods

- Learn a nonlinear function $f: \mathcal{R}^{d} \rightarrow \mathcal{R}$ from noisy samples

$$
\gamma_{i}=f\left(x_{i}\right)+\epsilon_{i} \text { for } i=1,2, \ldots n
$$

- $\epsilon_{i}$ are iid Gaussian noise with zero mean and variance $\lambda$
- Kernel Ridge Regression is a simple and yet powerful solution
- If $f(\cdot)$ is a GP with covariance $k: \mathcal{R}^{d} \times \mathcal{R}^{d} \rightarrow \mathcal{R}$, then the optimal estimator is,

$$
\begin{gathered}
\tilde{f}(x)=\sum_{i=1}^{n} \alpha_{i} k\left(x, x_{i}\right) \\
\alpha=\operatorname{argmin}_{\beta \in \mathbb{R}^{n}}\|\boldsymbol{K} \beta-\gamma\|_{2}^{2}+\lambda \beta^{\top} \boldsymbol{K} \beta
\end{gathered}
$$



## Kernel Method

- Kernel methods are expensive

- Computing all kernel entries takes $n \cdot n n z(X)+n^{2}$ time
- Even writing it down takes $n^{2}$ time and memory
- A single iteration of a linear system solver takes $n^{2}$ time
- For $n=100000, K$ has 10 billion entries. Takes 80 GB of storage


## Classical Solution: Dimensionality Reduction



- Storing $Z$ uses $O(n s)$ space and computing $Z Z^{\top} \alpha$ takes $O(n s)$ time.
- Orthogonalization, eigen-decomposition, and pseudo-inversion of $Z Z^{\top}$ all take just $O\left(n s^{2}\right)$ time.


## Efficient Low-Rank Approximation?

- Direct eigen decomposition, or even approximation via Krylov subspace methods are out of question since they at least require fully forming $K$


## Efficient Low-Rank Approximation?

- Direct eigen decomposition, or even approximation via Krylov subspace methods are out of question since they at least require fully forming $K$
- Sketching: a powerful approach to speeding up matrix problems
- Our approach: design a sketching solution for kernel low-rank approximation


## Feature Space Mapping

- Any positive definite kernel $k: \mathcal{R}^{d} \times \mathcal{R}^{d} \rightarrow \mathcal{R}$ defines a lifting $\varphi: \mathcal{R}^{d} \rightarrow \mathcal{R}^{D}$ such that for all $x, y \in \mathcal{R}^{d}$

$$
k(x, y)=\varphi(x)^{\top} \varphi(y)
$$

- The kernel computes the inner product between the lifted data points


## Feature Space Mapping

- Any positive definite kernel $k: \mathcal{R}^{d} \times \mathcal{R}^{d} \rightarrow \mathcal{R}$ defines a lifting $\varphi: \mathcal{R}^{d} \rightarrow \mathcal{R}^{D}$ such that for all $x, y \in \mathcal{R}^{d}$

$$
k(x, y)=\varphi(x)^{\top} \varphi(y)
$$

- The kernel computes the inner product between the lifted data points

$$
K=\boldsymbol{\phi}^{\top} \boldsymbol{\phi},
$$

where $\boldsymbol{\phi}$ is a $D \times n$ matrix whose $i^{\text {th }}$ column is the projection of $x_{i}$ into the feature space $\varphi\left(x_{i}\right)$

## Sketching the Feature Space

- Sketch the feature space

$$
K=\boldsymbol{\phi}^{\top} \boldsymbol{\phi} \approx \boldsymbol{\phi}^{\top} \Pi^{\top} \Pi \boldsymbol{\phi}
$$

## Sketching the Feature Space

- Sketch the feature space

$$
K=\boldsymbol{\phi}^{\top} \boldsymbol{\phi} \approx \boldsymbol{\phi}^{\top} \Pi^{\top} \Pi \boldsymbol{\phi}
$$

- Challenge: forming the feature matrix $\boldsymbol{\phi}$ explicitly is expensive as the feature space is typically high-dimensional (even infinite-dimensional)


## Sketching the Feature Space

- Sketch the feature space

$$
K=\boldsymbol{\phi}^{\top} \boldsymbol{\phi} \approx \boldsymbol{\phi}^{\top} \Pi^{\top} \Pi \boldsymbol{\phi}
$$

- Challenge: forming the feature matrix $\boldsymbol{\phi}$ explicitly is expensive as the feature space is typically high-dimensional (even infinite-dimensional)
- Goal: Design a sketch matrix $\Pi \in \mathcal{R}^{s \times D}$ such that $\Pi \cdot \varphi(x)$ is computable without needing to explicitly form $\varphi(x)$


## Kernel Sketching techniques

- The most popular method for kernel sketching is the Fourier Features Method of Rahimi \& Recht (Test of Time Award winner at NeurIPS'17)
- Works for shift invariant kernels, such as Gaussian kernel

$$
\varphi(x)_{\xi}=e^{-2 \pi i \xi^{\top} x} \text { for } \xi \in \mathcal{R}^{d}
$$

- $\Pi$ : Sampling matrix that samples frequencies $\xi$ from the $\operatorname{pdf} \hat{k}(\xi)$


## Kernel Sketching techniques

- The most popular method for kernel sketching is the Fourier Features Method of Rahimi \& Recht (Test of Time Award winner at NeurIPS'17)
- Works for shift invariant kernels, such as Gaussian kernel

$$
\varphi(x)_{\xi}=e^{-2 \pi i \xi^{\top} x} \text { for } \xi \in \mathcal{R}^{d}
$$

- П: Sampling matrix that samples frequencies $\xi$ from the $\operatorname{pdf} \hat{k}(\xi)$
- Avron, Kapralov, Musco, Musco, Velingker, Z.' 17: Tight bounds to get spectral approximation guarantee + Modified Fourier Sampling with optimal number of samples for Gaussian kernel in constant dimension


## Kernel Sketching techniques

- The most popular method for kernel sketching is the Fourier Features Method of Rahimi \& Recht (Test of Time Award winner at NeurIPS'17)
- Works for shift invariant kernels, such as Gaussian kernel

$$
\varphi(x)_{\xi}=e^{-2 \pi i \xi^{\top} x} \text { for } \xi \in \mathcal{R}^{d}
$$

- П: Sampling matrix that samples frequencies $\xi$ from the $\operatorname{pdf} \hat{k}(\xi)$
- Avron, Kapralov, Musco, Musco, Velingker, Z.' 17: Tight bounds to get spectral approximation guarantee + Modified Fourier Sampling with optimal number of samples for Gaussian kernel in constant dimension
- Avron, Kapralov, Musco, Musco, Velingker, Z.' 19: Optimal sampling strategy for Sinc kernel in dimension 1


## Kernel Sketching techniques

- The most popular method for kernel sketching is the Fourier Features Method of Rahimi \& Recht (Test of Time Award winner at NeurIPS'17)
- Works for shift invariant kernels, such as Gaussian kernel

$$
\varphi(x)_{\xi}=e^{-2 \pi i \xi^{\top} x} \text { for } \xi \in \mathcal{R}^{d}
$$

- П: Samf Works only for shift invariant kernels and $\operatorname{pdf} \hat{k}(\xi)$
- Avron, Ka constant dimensional datasets $s$ to get
spectral optimal number of samples for Gaussian kernel in constant dimension
- Avron, Kapralov, Musco, Musco, Velingker, Z.' 19: Optimal sampling strategy for Sinc kernel in dimension 1


## Polynomial Kernel

- In this work we focus on the important case of Polynomial Kernel

$$
k(x, y)=\left(x^{\top} y\right)^{q}
$$

- The lifting function for this kernel is $\varphi(x)=x^{\otimes q}$, where $\varphi(x) \in \mathcal{R}^{d^{q}}$ is defined as $\varphi(x)_{i_{1}, i_{2}, \cdots i_{q}}=x_{i_{1}} x_{i_{2}} \cdots x_{i_{q}}$ for $i_{1}, i_{2}, \cdots i_{q} \in\{1,2, \cdots d\}$


## Polynomial Kernel

- In this work we focus on the important case of Polynomial Kernel

$$
k(x, y)=\left(x^{\top} y\right)^{q}
$$

- The lifting function for this kernel is $\varphi(x)=x^{\otimes q}$, where $\varphi(x) \in \mathcal{R}^{d^{q}}$ is defined as $\varphi(x)_{i_{1}, i_{2}, \cdots i_{q}}=x_{i_{1}} x_{i_{2}} \cdots x_{i_{q}}$ for $i_{1}, i_{2}, \cdots i_{q} \in\{1,2, \cdots d\}$
- Goal: design a linear sketch $\Pi \in \mathcal{R}^{s \times d^{q}}$ such that $\Pi x^{\otimes q}$ is efficiently computable without needing to form $x^{\otimes q}$ explicitly


## Key Properties of Sketch

- Approximate Matrix Product: for every matrices $A, B \in \mathcal{R}^{d^{q} \times n}$ whp

$$
\left\|A^{\top} \Pi^{\top} \Pi B-A^{\top} B\right\|_{F} \leq \epsilon\|A\|_{F}\|B\|_{F}
$$

- Oblivious Subspace Embedding: for every $\lambda>0$ and every matrix $A \in \mathcal{R}^{d^{q} \times n}$ whp

$$
\frac{A^{\top} A+\lambda I}{1+\epsilon} \preccurlyeq A^{\top} \Pi^{\top} \Pi A+\lambda I \preccurlyeq \frac{A^{\top} A+\lambda I}{1-\epsilon}
$$

- Want: target dimension at most statistical dimension $\operatorname{tr}\left(A^{\top} A\left(A^{\top} A+\lambda I\right)^{-1}\right)$


## Prior Work: TensorSketch

- Originally introduced by Pagh and Pham [KDD 2013], [TOCT 2013]


## Prior Work: TensorSketch

- Originally introduced by Pagh and Pham [KDD 2013], [TOCT 2013]
- Avron, Nguyen, and Woodruff [NeurIPS 2014] proved:

1. Satisfies Approximate Matrix Product with probability $9 / 10$ if target dimension $s=\Omega\left(\frac{3{ }^{\mathrm{q}}}{\epsilon^{2}}\right)$
2. Satisfies Oblivious Subspace Embedding with probability 9/10 if target dimension $s=\Omega\left(\frac{3^{\mathrm{q}}}{\epsilon^{2}} \cdot s_{\lambda}^{2}\right)$
3. Time to sketch the tensor $x^{\otimes q}$ is $\tilde{O}(\mathrm{q} s+\mathrm{q} \cdot n n z(x))$

$$
\text { Statistical Dimension } s_{\lambda}:=\boldsymbol{\operatorname { t r }}\left(K(K+\lambda I)^{-1}\right)
$$

## Prior Work: TensorSketch

- Originally introduced by Pagh and Pham [KDD 2013], [TOCT 2013]
- Avron, Nguyen, and Woodruff [Neu

1. Satisfies Approximate Matrix Pr dimension $s=\Omega\left(\frac{3^{q}}{\epsilon^{2}}\right)$
2. Satisfies Oblivious Subspace Em

Main contribution: improve the exponential target
dependence on q to polynomial
target dimension $s=\Omega\left(\frac{3^{9}}{\epsilon^{2}} \cdot s_{\lambda}^{2}\right)$
3. Time to sketch the tensor $x^{\otimes q}$ is $\tilde{O}(\mathrm{q} s+\mathrm{q} \cdot n n z(x))$

$$
\text { Statistical Dimension } s_{\lambda}:=\boldsymbol{\operatorname { t r }}\left(K(K+\lambda I)^{-1}\right)
$$

## Prior Work: TensorSketch

- Originally introduced by Pagh and Pham [KDD 2013], [TOCT 2013]
- Avron, Nguyen, and Woodruff [Neurine ani nl manuad.

1. Satisfies Approximate Matrix Pro dimension $s=\Omega\left(\frac{3{ }^{\mathrm{q}}}{\epsilon^{2}}\right)$

Contribution 2: improve the quadratic dependence on $s_{\lambda}$ to linear
2. Satisfies Oblivious Subspace Embeuums vitir provavinty JIfo if target dimension $s=\Omega\left(\frac{3^{q}}{\epsilon^{2}} \cdot s_{\lambda}^{2}\right)$
3. Time to sketch the tensor $x^{\otimes q}$ is $\tilde{O}(\mathrm{q} s+\mathrm{q} \cdot n n z(x))$

$$
\text { Statistical Dimension } s_{\lambda}:=\boldsymbol{\operatorname { t r }}\left(K(K+\lambda I)^{-1}\right)
$$

## Prior Work: TensorSketch

- Originally introduced by Pagh and Pham [KDD 2013], [TOCT 2013]
- Avron, Nguyen, and Woodruff [NeurIPS 2014] proved:

1. Satisf Contribution 3: improve uct with probability 9/10 if targe the success probability to
2. Satis
$1-\frac{1}{\operatorname{poly}(n)} \quad$ dding with probability $9 / 10$ if target dimension $s=\Omega\left(\overline{\epsilon^{2}} \cdot s_{\lambda}^{2}\right)$
3. Time to sketch the tensor $x^{\otimes q}$ is $\tilde{O}(\mathrm{q} s+\mathrm{q} \cdot n n z(x))$

$$
\text { Statistical Dimension } s_{\lambda}:=\boldsymbol{\operatorname { t r }}\left(K(K+\lambda I)^{-1}\right)
$$

## Main Results

- Theorem 1: there exists a distribution on linear sketches $\Pi \in \mathcal{R}^{s \times d^{q}}$ such that:

1. If target dimension $s=\Omega\left(\frac{q}{\epsilon^{2}}\right)$ then $\Pi$ has the Approximate Matrix Product property with probability 9/10
2. If target dimension $s=\Omega\left(\frac{q}{\epsilon^{2}} \cdot s_{\lambda}^{2}\right)$ then $\Pi$ is an Oblivious Subspace Embedding with probability 9/10
3. For any $x \in \mathcal{R}^{d}, \Pi \cdot x^{\otimes q}$ is computable in time $\tilde{O}(\mathrm{q} s+\mathrm{q} \cdot n n z(x))$

## Main Results

- Theorem 2: there exists a distribution on linear sketches $\Pi \in \mathcal{R}^{s \times d^{q}}$ such that:

1. If target dimension $s=\widetilde{\Omega}\left(\frac{q^{4}}{\epsilon^{2}}\right)$ then $\Pi$ has the Approximate Matrix Product property with high probability
2. If target dimension $s=\widetilde{\Omega}\left(\frac{q^{4}}{\epsilon^{2}} \cdot s_{\lambda}\right)$ then $\Pi$ is an Oblivious Subspace Embedding with high probability
3. For any vector $x \in \mathcal{R}^{d}$, the product $\Pi \cdot x^{\otimes q}$ is computable in time $\tilde{o}\left(q s+q^{5} \epsilon^{-2} n n z(x)\right)$

## Review: TensorSketch

$$
\Pi x^{q}=\mathcal{F}^{-1}\left[\left(\mathcal{F} C_{1} x\right) \circ\left(\mathcal{F} C_{2} x\right) \circ \cdots \circ\left(\mathcal{F} C_{q} x\right)\right]
$$

- $\mathcal{F}$ is the Fourier transform matrix and $C_{1}, C_{2}, \cdots C_{q} \in \mathcal{R}^{s \times d}$ are independent copies of CountSketch


## Review: TensorSketch

$$
\Pi x^{q}=\mathcal{F}^{-1}\left[\left(\mathcal{F} C_{1} x\right) \circ\left(\mathcal{F} C_{2} x\right) \circ \cdots \circ\left(\mathcal{F} C_{q} x\right)\right]
$$

- $\mathcal{F}$ is the Fourier transform matrix and $C_{1}, C_{2}, \cdots C_{q} \in \mathcal{R}^{s \times d}$ are independent copies of CountSketch
- The second moment of this estimator for $x=\{1\}^{d}$

$$
\mathbb{E}\left[\left\|\Pi x^{\otimes q}\right\|_{2}^{4}\right] \geq \frac{3^{q}}{2 s^{2}}\left\|x^{\otimes q}\right\|_{2}^{4}
$$

## Our Sketch Construction

- Every node is an independent instance of some base sketch
- Leaves: sketch the input vector
- Internal nodes: sketch the tensor product of their children



## Our Sketch Construction

- Every leaf is a sketch that runs in input sparsity time
- Internal nodes support fast application time



## Our Sketch Construction

- Intermediate nodes tensor only twice

- In particular, there is no exponential dependence on q
$s$ : target dimension of
intermediate sketches


## OSE for Gaussian Kernel

- The polynomial dependence of our sketch on the degree q leads to significant improvements on sketching the Gaussian kernel in high-d


## OSE for Gaussian Kernel

- The polynomial dependence of our sketch on the degree q leads to significant improvements on sketching the Gaussian kernel in high-d


## Prior work

- Fast multipole method of Greengard and Rokhlin: suffers from curse of dimensionality $(\log n)^{d}$
- Fourier features method of Rahimi \& Recht: significantly suboptimal runtime of $\frac{n}{\lambda} \cdot n n z(X)$
- Modified Fourier features of Avron, Kapralov, Musco, Musco, Velingker, $\mathbf{Z}^{\prime}$ 17: Optimal for constant dimensions d but does not apply to high dimensional data


## OSE for Gaussian Kernel

- Theorem 3: for any dataset $x_{1}, x_{2}, \cdots x_{n} \in \mathbb{R}^{d}$ such that $\left\|x_{i}\right\|_{2}^{2} \leq r$ if $K \in \mathbb{R}^{n \times n}$ is the Gaussian kernel matrix defined as $K_{i, j}=e^{-\left\|x_{i}-x_{j}\right\|_{2}^{2}}$ there exists an algorithm that computes $Z \in \mathbb{R}^{n \times s}$ such that:

1. If target dimension $s=\widetilde{\Omega}\left(\frac{r^{5}}{\epsilon^{2}} \cdot s_{\lambda}\right)$ then $\mathrm{ZZ}^{\top}$ is an Oblivious Subspace Embedding for kernel $K$ with high probability
2. The runtime to compute $Z$ is $\tilde{O}\left(r^{6} \epsilon^{-2} n s_{\lambda}+r^{6} \epsilon^{-2} n n z(X)\right)$
